

On Conway's Numbers and Games

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UNIVERSITY
OF AMSTERDAM



INSTITUTE FOR LOGIC,
LANGUAGE AND COMPUTATION

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Dedekind cuts

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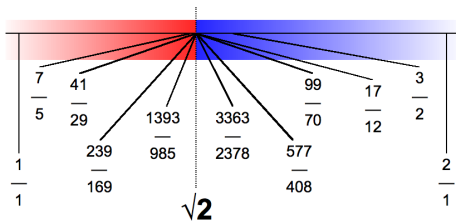
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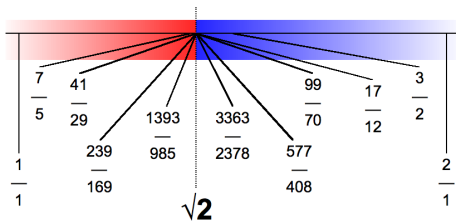


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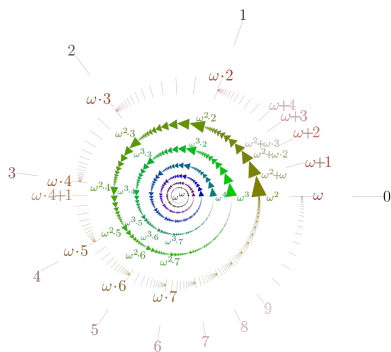
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Conway's Numbers

Construction

If L, R are two sets of numbers, and for all $l \in L, r \in R$, we have $l \not\geq r$, then $\{L \mid R\}$ is a number.

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Convention

If $x = \{L \mid R\}$, we write x^L for the typical member of L , x^R for the typical member of R , and $\{x^L \mid x^R\}$ for x itself.

We write No for the (proper) class of numbers.

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Definitions

We say $x \geq y$ iff $x^R \not\leq y$ and $x \not\leq y^L$ (for all x^R and y^L).

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$$x + y = \{x + y^L, x^L + y \mid x + y^R, x^R + y\}$$

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Finally, we let

$$xy = \{x^L y + xy^L - x^L y^L, x^R y + xy^R - x^R y^R \mid x^L y + xy^R - x^L y^R, x^R y + xy^L + x^R y^L\}$$

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$$-2 = \{ | -1 \} = \{ | -1, 0 \} = \{ | -1, 1 \} = \{ | -1, 0, 1 \}$$

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$$-\frac{1}{2} = \{-1 | 0\} = \{-1 | 0, 1\}$$

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Prove that $1 + 1 = 2$ and that $\frac{1}{2} + \frac{1}{2} = 1$.

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What is the number $\{0 \mid \frac{1}{2}\}$? What about $\{\frac{1}{2} \mid 1\}$? $\{1 \mid 2\}$? $\{2 \mid\}$?

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We also 'recreate' dyadic rationals. For example,

$\{\text{dyadic rationals} < \frac{3}{8} \mid \text{dyadic rationals} > \frac{3}{8}\}$ turns out to be $\frac{3}{8}$.

Some more Numbers

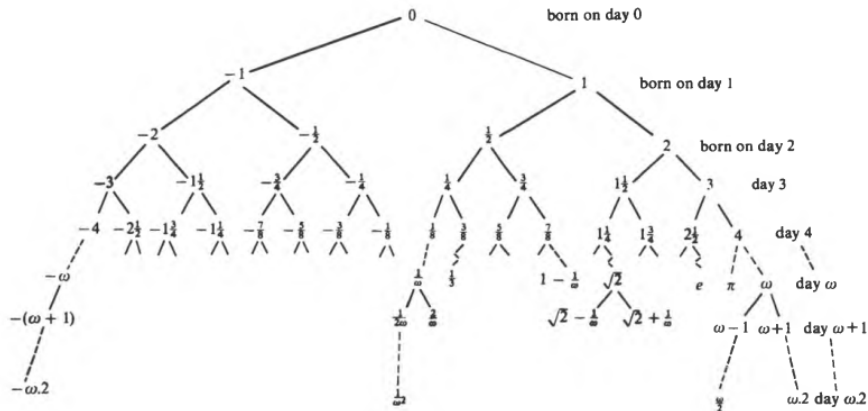


FIG. 0. When the first few numbers were born.

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Remark

Multiplication (of games) preserves identity but not equality.

Properties of order and equality

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Theorem

For any numbers x and y , we have $x^L < x < x^R$, and $x \geq y$ or $x \leq y$.

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Summary

Numbers are totally ordered by \geq .

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For all x, y, z , we have $x + 0 = x$, $x + y \equiv y + x$ and $(x + y) + z \equiv x + (y + z)$.

Group properties

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For all x, y, z , we have $x + 0 = x$, $x + y \equiv y + x$ and $(x + y) + z \equiv x + (y + z)$.

Theorem (Properties of negation)

For all x, y , we have $-(x + y) \equiv -x + -y$,

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Proof.

Induction,

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Induction, left as an exercise to the audience. □

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Proof.

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Summary

The operations $+$, $-$ and 0 induce a Group structure on Games.

Ordered group properties

Theorem

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Theorem

For all x, y, z , we have $y \geq z$ iff $x + y \geq x + z$.

Corollary

If $x_1 = x_2$, then $x_1 + y = x_2 + y$.

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Theorem

For all x, y, z , we have $y \geq z$ iff $x + y \geq x + z$.

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If $x_1 = x_2$, then $x_1 + y = x_2 + y$.

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Summary

Numbers form a totally ordered Group.

Properties of multiplication

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Multiplication induces a Ring structure on Numbers and on Games.

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Multiplication induces a Ring structure on Numbers and on Games.

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It is also possible to define division, square roots, etc, turning the Class of numbers into totally ordered Field with many properties.

How does one single out the real numbers from the Class of numbers?

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Definition

A number x is a real number iff $-n < x < n$ for some integer n and

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- *Dyadic rationals are real numbers.*
- *Real numbers are closed under the field operations.*

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A game α is an ordinal iff it has an expression of the form $\alpha = \{L \mid \}$.

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They correspond to the *natural* sum and product, obtained by treating the Cantor normal form of an ordinal as a polynomial.

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Remark

In \mathbb{N}_0 , one can define sums indexed over any ordinal. Analytic functions (such as \log , \sin , \exp) can be defined as power series. One could also use various small subfields of \mathbb{N}_0 as a model for non-standard analysis.

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Theorem (Euclidean division)

If a and b are integers with b positive, there are unique integers q and r such that $a = bq + r$ and $0 \leq r < b$.

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Almost every number-theoretical problem can be rephrased so as to yield a new problem in \mathbb{O}_Z , so we get a jackdaw's nest of problems of various kind. Examples include Waring's problem, continued fractions, Pellian equations.

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Our players Left and Right are unwilling to play games that may go on forever (they are both busy people, with heavy political responsibilities).

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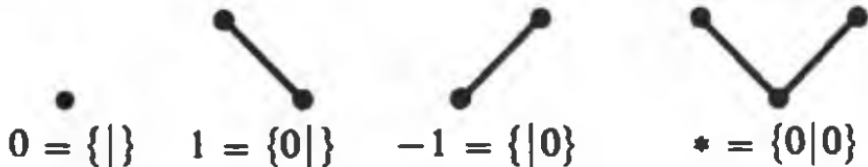


FIG. 4. The simplest games.

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Theorem

Each game belongs to one of the outcome classes above.

Definition (Negation of a game)

If we reverse the roles of Left and Right in a game G , we obtain the game

$$-G = \{-G^R \mid -G^L\}.$$

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How to treat fuzzy numbers

Some games are not numbers:

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Definition

A *short game* is one which has only finitely many positions.

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Theorem

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Conclusion

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There has been some work on Surreal numbers, mainly by Donald Knuth, Harry Gonshor, Norman Alling, Philip Ehrlich and Martin Kruskal. This work is mainly focused on nonstandard analysis.

There has also been some work on Numbers from an algebraic perspective, proving universal properties of the Field / Group of numbers.